A combinatorial approach to Macdonald *q*, *t*-symmetry via the Carlitz bijection

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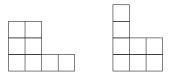
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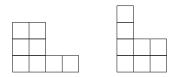
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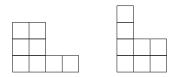


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- Related to classical Macdonald polynomials P_λ by a transformation, arise naturally in the geometry of the Hilbert scheme of points in the plane. (Haiman)
- q, t-symmetry (via geometry): $\widetilde{H}_{\mu}(X; q, t) = \widetilde{H}_{\mu^*}(X; t, q)$

Combinatorial Formula:

$$\widetilde{H}_{\mu}(X;q,t) = \sum_{\sigma: \mu o \mathbb{Z}_+} q^{\mathsf{inv}(\sigma)} t^{\mathsf{maj}(\sigma)} x^{\sigma}$$

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$$\mu = (3, 2, 2), \ \sigma = \boxed{\begin{array}{c|c} 5 & 1 \\ 2 & 1 \\ 2 & 3 & 2 \end{array}}$$

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- inv and maj are statistics on fillings that generalize inv and maj on permutations.
- q, t-symmetry not obvious from this formula

• Let $\pi_1 \pi_2 \cdots \pi_n$ be a permutation of [n]. Then

$$\mathsf{inv}(\pi) = |\{(i,j): i < j, \pi_i > \pi_j\}|,$$
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inv and maj are equidistributed:

$$\sum_{\pi\in S_n} q^{\mathsf{inv}(\pi)} = \sum_{\pi\in S_n} q^{\mathsf{maj}(\pi)} = (1)(1\!+\!q)(1\!+\!q\!+\!q^2)\cdots(1\!+\!q\!+\!\cdots\!+\!q^{n-1})$$

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• Combinatorial proofs: Find a "weight-preserving" bijection $\phi: S_n \to S_n$, i.e. a bijection such that $\operatorname{maj}(\phi(\pi)) = \operatorname{inv}(\pi)$. Several such maps have been found (Carlitz, Foata,...)

▶ **Carlitz codes:** Let $C_n = \{c_1c_2 \cdots c_n : \forall i, 0 \le c_i \le n-i\}$. Define the *weight* of a code $c \in C_n$ to be $|c| = \sum_i c_i$. Then

$$\sum_{c \in \mathcal{C}_n} q^{|c|} = (1)(1+q)(1+q+q^2) \cdots (1+q+\cdots+q^{n-1})$$

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► Carlitz bijection: Composite φ : S_n → S_n of two weight-preserving bijections

$$S_n \xrightarrow{\text{majcode}} C_n \xleftarrow{\text{invcode}} S_n.$$

Weight-preserving: $|majcode(\pi)| = maj(\pi)$ and $|invcode(\pi)| = inv(\pi)$.

$$S_n \xrightarrow{\text{majcode}} C_n \xleftarrow{\text{invcode}} S_n$$

majcode: Remove entries starting with the largest, c_i records the amount maj decreases at the *i*th step:

Word	maj	Ci
51423	4	

majcode(51423) =

$$S_n \xrightarrow{\text{majcode}} C_n \xleftarrow{\text{invcode}} S_n$$

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Word	maj	Ci
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$$majcode(51423) = 2$$

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Word	maj	Ci
51423	4	
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123		

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51423	4	
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Word	maj	Ci
51423	4	
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12		

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51423	4	
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$$51423 \rightarrow 34125.$$

 Can extend the Carlitz bijection from permutations to words of any content.

Conjugate Symmetry in q and t

Recall that

$$\widetilde{H}_{\mu}(X;q,t) = \widetilde{H}_{\mu^*}(X;t,q).$$

• Take the coefficient of $x^{\alpha} = x_1^{\alpha_1} x_2^{\alpha_2} \cdots$ on both sides:

$$\sum_{\substack{\sigma:\mu o\mathbb{Z}_+\ |\sigma^{-1}(i)|=lpha_i}} q^{\mathsf{inv}(\sigma)} t^{\mathsf{maj}(\sigma)} = \sum_{\substack{
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Combinatorial proof: Need a bijection from fillings of µ to fillings of µ* that preserves content and switches inv and maj simultaneously.

$$\sigma = \begin{bmatrix} 6 & 3 \\ 1 & 5 & 7 \\ 2 & 4 & 1 \end{bmatrix}$$

maj is the sum of the maj's of the columns (top to bottom).

• $maj(\sigma) = maj(612) + maj(354) + maj(71) = 1 + 2 + 1 = 4$

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- Inversions: A pair (a, b) with a to the left of b in a row is usually an inversion if and only if a > b, except if the entry c directly below a has value between a and b then the opposite is true.

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- ▶ Bottom row: (2,1)

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- Second row: (1,5)



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- Bottom row: (2,1), (4,1)
- Second row: (1,5), (1,7)
- Top row: (6,3)

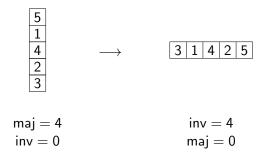
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- Bottom row: (2,1), (4,1)
- Second row: (1,5), (1,7)
- ▶ Top row: (6,3)

•
$$inv(\sigma) = 5$$

One-row shapes reduce to words

- If μ = (n) and σ is a filling of μ, then maj(σ) = 0 and inv(σ) = inv(w(σ)) where w(σ) is the reading word of σ.
- Similarly if ρ fills μ^{*} then inv(ρ) = 0 and maj(ρ) = maj(w(σ)).



Result: Hook shapes (G.)

Given a filling σ of a hook shape, define invcode and majcode according to the invcode and majcode of the row and column respectively.

Leftmost 0 of invcode matches rightmost 0 of majcode.

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Given a filling σ of a hook shape, define invcode and majcode according to the invcode and majcode of the row and column respectively.

- Leftmost 0 of invcode matches rightmost 0 of majcode.
- Now interchange and reverse the two codes!

Hall-Littlewood specialization: q = 0

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- ► Hall-Littlewood polynomials: $\widetilde{H}_{\mu}(X; t) = \widetilde{H}_{\mu}(X; 0, t)$
- Symmetry problem restricts to fillings of µ with inv = 0 and fillings of µ^{*} with maj = 0. Need generalized Carlitz codes.

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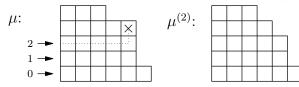
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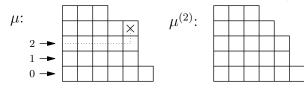
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• Define $C_{\mu} = \{c_1 \cdots c_n : z_n^{c_1} \cdots z_1^{c_n} \in \mathcal{B}_{\mu}\}.$

Generalizing Carlitz

Want weight-preserving bijections

$$\mathcal{F}_{\mu}|_{\mathsf{inv}=0} \xrightarrow{\mathsf{majcode}} \mathcal{C}_{\mu} \xleftarrow{\mathsf{invcode}} \mathcal{F}_{\mu^*}|_{\mathsf{maj}=0}$$

where

$$\mathcal{F}_{\mu}|_{\mathsf{inv}=\mathsf{0}} = \{\sigma: \mu \to \mathbb{Z}_+ \mid \mathsf{inv}(\sigma) = \mathsf{0}\}$$

and

$$\mathcal{F}_{\mu^*}|_{\mathsf{maj}=\mathsf{0}} = \{\rho: \mu^* \to \mathbb{Z}_+ \mid \mathsf{maj}(\rho) = \mathsf{0}\}.$$

Here the *weight* of a code $c \in C_{\mu}$ is $|c| = \sum c_i$, so the maj statistic on the left and the inv statistic on the right will be sent to this weight statistic on C_{μ} .

Let ρ be a filling of μ* having maj(ρ) = 0. Order its entries by size with ties broken in reading order, forming a totally ordered alphabet A = {a₁,..., a_n}.

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- Define invcode(ρ) = c₁c₂ ··· c_n where c_i is the number of attacking pairs having a_i as its smaller entry.

$$\rho = \frac{\begin{vmatrix} 3 & 2 \\ 4 & 6 & 1 \\ 5 & 6 & 2 \end{vmatrix}$$

$$A = 12234566$$
invcode(ρ) = 21200100

Theorem (G.)

The map invcode is a weight-preserving bijection

$$\mathsf{invcode}: \mathcal{F}_{\mu^*}|_{\mathsf{maj}=0} \xrightarrow{\sim} \mathcal{C}_{\mu}.$$

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$$\mathcal{F}_{\mu}|_{\mathsf{inv}=0} \xrightarrow{\mathsf{majcode}} \mathcal{C}_{\mu} \xleftarrow{\mathsf{invcode}} \mathcal{F}_{\mu^*}|_{\mathsf{maj}=0}$$

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- To extend majcode to fillings, want to remove the largest entry and record how much the maj decreases by at each step.
- How to remove largest entry from an inversion-free filling σ ?

5		
7	1	2
3	4	6

• The recursion defining \mathcal{B}_{μ} implies that we should remove the largest entry from σ so that, if the major index decreases by d, the resulting shape is $\mu^{(d)}$.

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- Via a different approach, can construct a map majcode for three-row shapes with general entries.

Results

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The Carlitz bijection can be extended to prove q, t-symmetry in the following cases:

• Hook shapes: When μ is a hook shape.

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- Hook shapes: When μ is a hook shape.
- q = 0, distinct entries: For fillings with distinct entries (and any shape μ) when one of the statistics is zero.
- ▶ q = 0, $\ell(\mu) \leq 3$: When one of the statistics is zero and μ has at most three parts, showing that

$$\widetilde{H}_{\mu}(X;0,t)=\widetilde{H}_{\mu^*}(X;t,0)$$

for such shapes.

Acknowledgments

 Thanks to Mark Haiman and Angela Hicks for their help and feedback.

THANK YOU!