

Designs

Problem in agriculture: (Experimental designs)

- v different varieties of grain, want to compare them

Grain types: 1, 2, 3, 4, 5, 6, 7

- Also want to vary conditions - soil, temperature, etc, so you make b blocks, each of the same size k, in terms of number of grains tested, in different conditions

$k=3$: $\{1, 2, 4\}$, $\{2, 3, 5\}$, ...
soil A soil B

$b=7$ (7 types of soil)

- Too inefficient to test every grain in every soil, but want to compare all pairs of grains to determine the best one:
 \uparrow
 $t=2$

Soil A: 1 2 3

Soil B: 3 4 5

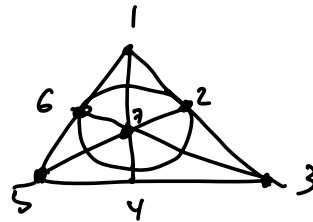
5 6 1

⋮ 1 4 7

⋮ 2 5 7

⋮ 3 6 7

2 4 6



"Fano plane"

7 blocks, each pair of grains appears together in exactly 1 block. Each grain is tested 3 times

- In general have each pair occur in λ blocks each. Each variety occurs in r blocks each.

Def: A t - (v, k, λ) design is a collection of blocks \mathcal{B} consisting of $\binom{v-t}{k-t}$ k -elt subsets of a set X of size v (elts of X are called points) s.t.
 * Every t -elt subset of X is in exactly λ blocks.

Ex: The Fano plane above is a 2 - $(7, 3, 1)$ design.

Def: For $\lambda=1$, called a Steiner system.

Note: We'll be assuming repeated blocks are not allowed.

Ex: $\binom{[n]}{k}$ is a t - (n, k, λ) design for what λ ? Ans: $\lambda = \binom{n-t}{k-t}$.

Lemma: Any t -design is also an s -design for any $s \leq t$.

Pf: Let (X, \mathcal{B}) be a t - (v, k, λ) design, and let $s \leq t$. Then each s -size subset $S \subseteq X$ is contained in exactly $\binom{v-s}{t-s}$ sets $T \subseteq X$ of size t . Each T occurs in λ blocks $B \in \mathcal{B}$, so this gives $\lambda \binom{v-s}{t-s}$ occurrences of S in blocks B , where we have overcounted by the number of T containing S within any given block B containing S , which is $\binom{k-s}{t-s}$.

So (X, \mathcal{B}) is an s - (v, k, λ') design where $\lambda' = \lambda \frac{\binom{v-s}{t-s}}{\binom{k-s}{t-s}}$.

Lemma: Let r be the number of blocks each $x \in X$ occurs in, in a design. Then

$$bk = vr$$

"Burger King = Virtual Reality"

Pf: $b \cdot k = \# \text{ pairs } (x, B): x \in B, B \in \mathcal{B}$
 $= v \cdot r$

(note: r exists by previous lemma)

Lemma: In a 2-design, have $r(k-1) = (v-1)\lambda$.

Pf: $r = \# \text{ blocks } x \in X \text{ is contained in}$
 $k-1 = \# \text{ ways to pick another elt } y \text{ to pair w/ } x \text{ in a block } B$
 $v-1 = \# \text{ elts } y \text{ besides } x \text{ in } X$
 $\lambda = \# \text{ blocks containing } (x, y)$

Fix $x \in X$. $(v-1)\lambda = \# \text{ pairs } (y, B) \text{ s.t. } x, y \in B.$
 $= r(k-1).$

Def: Incidence matrix of a design: Write each block as a 0-1 row vector.

Ex: Fano plane is

	1	2	3	4	5	6	7
(1	1	1				
			1	1	1		
	1				1	1	
	1			1			1
		1			1	1	
			1			1	1
		1		1	1		

Can rearrange rows:

$$\begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 \end{pmatrix}$$

↙ symmetric matrix!

Def: The dual of a design $\mathcal{D} = (X, \mathcal{B})$ is written \mathcal{D}^T and consists of transposing the incidence matrix, i.e.

$$X^T = \mathcal{B}$$

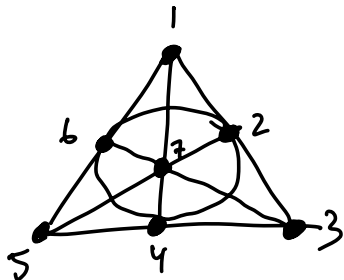
$$\mathcal{B}^T = \{ \{ B \ni x \} : x \in X \}$$

Symmetric design if $\mathcal{D}^T \cong \mathcal{D}$.

Isomorphism: $(X, \mathcal{B}) \rightarrow (X', \mathcal{B}')$ is a bijection $X \rightarrow X'$ that induces a bijection $\mathcal{B} \rightarrow \mathcal{B}'$.

Note: There are nontrivial automorphisms - isomorphisms to itself.

Ex: Automorphism of the Fano plane:



$$6 \leftrightarrow 2$$

$$5 \leftrightarrow 3$$

Ex: A $2-(9, 3, 1)$ design:

123
456
789
147
258
369
168
249
357
159
267
348

1 2 3
: ' '
4 5 6
: ' '
7 8 9

Dual is not a 2-design, but is a 1-design,

The dual of a 1-design is a 1-design

Above ex is a $1-(9, 3, 4)$ design

Its dual is a $1-(12, 4, 3)$ design.

In fact, 2-designs' duals are 2-designs
iff their incidence matrix is square, i.e. $b=v$:

Thm (Fisher's Inequality) In a 2-design with
 $k < v$, we have $b \geq v$.

Pf: First note that if M is the inc. matrix,

$$M = \begin{matrix} & x_1 & \dots & x_v \\ \begin{matrix} B_1 \\ \vdots \\ B_b \end{matrix} & \begin{pmatrix} & & & \\ & & & \\ & & & \\ & & & \end{pmatrix} \end{matrix}$$

$M^T M =$ matrix whose (i, j) entry is the dot product of column i w/ col j in M

$=$ matrix whose (i, j) entry is

$$\begin{cases} \# B \text{ containing } x_i \text{ and } x_j & \text{if } i \neq j \\ \# B \text{ containing } x_i & \text{if } i = j \end{cases}$$

$$= \begin{cases} \lambda & \text{if } i \neq j \\ r & \text{if } i = j \end{cases}$$

$$= (r - \lambda) I_v + \lambda J_v.$$

$$J_v = \begin{pmatrix} (1) & (1) & (1) & (1) \\ (1) & (1) & (1) & (1) \\ (1) & (1) & (1) & (1) \\ (1) & (1) & (1) & (1) \end{pmatrix}$$

Want to show $M^T M$ is nonsingular, which would show $b \geq v$. So want to compute

$$\det(M^T M) = \det((r - \lambda) I_v + \lambda J_v)$$

$=$ product of eigenvalues.

The all-1's vector $u = \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix}$ is an eigenvector

w/ eigenvalue $r - \lambda + \lambda v = r + \lambda(v-1) = r + r(k-1) = rk$.

The orthogonal vectors to u , like $\begin{pmatrix} 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \end{pmatrix}$,

are also eigens w/ eigenvalue $r - \lambda$

$\Rightarrow \det = rk(r - \lambda)^{v-1}$.

To check $r - \lambda \neq 0$: we have

$$r(k-1) = (v-1)\lambda$$

$$(r - \lambda)(k-1) = \lambda(v-1) - \lambda(k-1)$$

$$(r - \lambda)(k-1) = \lambda(v-k)$$

$\uparrow > 0 \quad \quad \quad \uparrow > 0$ by assumption

so this is > 0 .

Thus $\det M^T M \neq 0$, so $b \geq v$.

Thm: TFAE (and such a 2-design is called square) for

a 2-design w/ $k < v$:

- a) $b = v$
- b) $r = k$
- c) Any two blocks have λ common points
- d) Any two blocks have a constant number of common points

Pf: (a) \Leftrightarrow (b) by $bk = vr$

(b) \Rightarrow (c): If $r = k$, $t \leq n$ $M J_v = J_r M = k J_{r \times v}$

$$\begin{aligned} \text{So } M M^T &= M(M^T M)M^{-1} \\ &= M((r-\lambda)I_v + \lambda J_v)M^{-1} \\ &= ((r-\lambda)I_r + \lambda J_r)M M^{-1} \\ &= (r-\lambda)I_r + \lambda J_r \end{aligned}$$

and so any two blocks have λ pts in common.

(c) \Rightarrow (d) trivially

(d) \Rightarrow (a): If (d) is true then the dual design is a 2-design, and so $b \geq v$ and $v \geq b$ and so $b = v$.

Hadamard designs

Def: A Hadamard matrix is an $n \times n$ matrix with ± 1 entries s.t. $H H^T = I_n$.

i.e. each row \vec{r} has $\langle \vec{r}, \vec{r} \rangle = n$ and two rows $\vec{r}_1 \neq \vec{r}_2$ have $\langle \vec{r}_1, \vec{r}_2 \rangle = 0$.

Ex:

$$(1) \quad \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad \begin{pmatrix} -1 & 1 & 1 & 1 \\ 1 & -1 & 1 & 1 \\ 1 & 1 & -1 & 1 \\ 1 & 1 & 1 & -1 \end{pmatrix} \quad \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix}$$

3x3 Hadamard?

$$\begin{pmatrix} 1 & 1 & 1 \\ 1 & -1 & 1 \end{pmatrix} \quad \times$$

→ must have even size (after 1)

Why Hadamard matrices? Equality case in

"Hadamard det thm": if all entries of M have $|M_{ij}| \leq 1$,
 $|\det(M)| \leq n^{n/2}$.

Note: Can multiply any row or col by -1 and still have a Hadamard matrix. Also swap rows or columns. Normalized:

$$\begin{pmatrix} 1 & 1 & 1 & \dots & 1 \\ \vdots & & & & \vdots \\ \vdots & & & & \vdots \\ \vdots & & & & \vdots \\ \vdots & & & & \vdots \end{pmatrix}$$

Ex: Normalization of

$$\begin{pmatrix} -1 & 1 & 1 & 1 \\ 1 & -1 & 1 & 1 \\ 1 & 1 & -1 & 1 \\ 1 & 1 & 1 & -1 \end{pmatrix} \quad \text{is}$$

$$\begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & 1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix}$$

Lemma: Hadamard matrices must have $n = 1, 2, 4t$.

Pf. Suppose H Hadamard, ^{size > 3 ,} consider first 3 rows, normalize and arrange cols so they look like:

$$\begin{pmatrix} | & | & | & \dots & | & | & | & \dots & | & | & | & \dots & | & | & | & \dots & | \\ | & | & | & \dots & | & | & | & \dots & | & | & | & \dots & | & | & | & \dots & | \\ | & | & | & \dots & | & | & | & \dots & | & | & | & \dots & | & | & | & \dots & | \end{pmatrix}$$

$\underbrace{\hspace{10em}}_x \quad \underbrace{\hspace{10em}}_y \quad \underbrace{\hspace{10em}}_z \quad \underbrace{\hspace{10em}}_w$

Then $x + y = z + w$ and $x + z = y + w$ and
 $x + w = y + z$

So $x = y = z = w$.

Thus n is div. by 4.

Conjecture (OPEN!) There \exists a Hadamard matrix of order $4t$ for every t .

Connection to designs

Prop: Let H be a normalized Hadamard matrix (first row and first column all 1's).

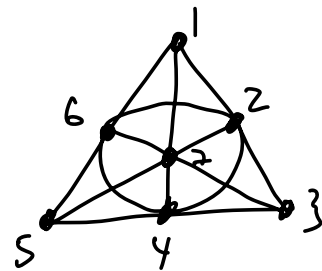
Let M be the matrix formed by deleting the first row, col from H and replacing every $+1$ by 1 and -1 by 0 .

Then M is the incidence matrix of a $2-(4t-1, 2t-1, t-1)$ design.

If we instead replace $+1$ by 0 and -1 by 1 , get a $2-(4t-1, 2t, t)$ design.

Ex:

$$\begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 \\ 1 & 1 & -1 & -1 & 1 & -1 & -1 & 1 \\ 1 & 1 & -1 & -1 & -1 & 1 & 1 & -1 \\ 1 & -1 & 1 & -1 & 1 & -1 & 1 & -1 \\ 1 & -1 & -1 & 1 & -1 & 1 & 1 & 1 \\ 1 & -1 & -1 & 1 & 1 & 1 & -1 & -1 \\ 1 & -1 & 1 & -1 & -1 & 1 & -1 & 1 \end{pmatrix}$$



↖ Hadamard matrix order 8
($t=2$)

Duality - an isom. between \mathcal{D} and its dual

(for a symmetric

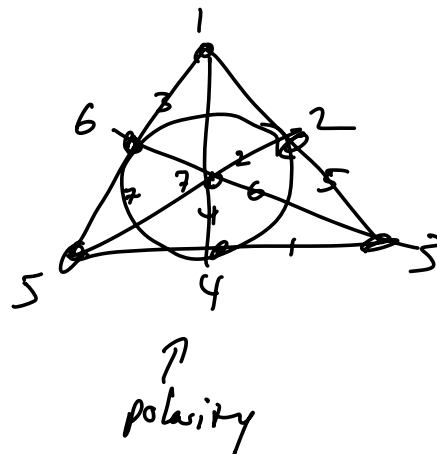
design \mathcal{D} with $\mathcal{D} \cong \mathcal{D}^T$):

a pair of bijections

$$\sigma: X \rightarrow \mathcal{B}$$


$$\tau: \mathcal{B} \rightarrow X$$

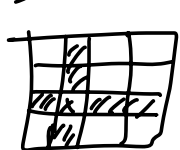
s.t. $x \in \mathcal{B}$ iff
 $\tau(\mathcal{B}) \in \sigma(x)$



If $\sigma\tau = \text{id}_{\mathcal{B}}$ and $\tau\sigma = \text{id}_X$,
 say the duality is a polarity.

Ex: A $2-(16, 6, 2)$ -design:

$X = \left\{ \begin{array}{l} \text{with} \\ \text{squares in: } \end{array} \right\}$ 

$\mathcal{B} = \left\{ \text{blocks of the form } \right\}$ 

(6 shaded squares in vertical
 or horiz alignment w/
 a given square x .)

Polarity: $x \rightarrow$ its associated block.

In this polarity, no x is in its own block.

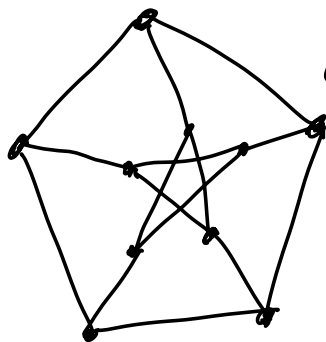
Strongly Regular Graphs (SRG's)

Regular graph - all vertices the same degree

Def. An SRG w/ parameters (n, k, λ, μ) is a graph w/ n vertices, not K_n or its complement, s.t. # common neighbors of any two vertices x, y is

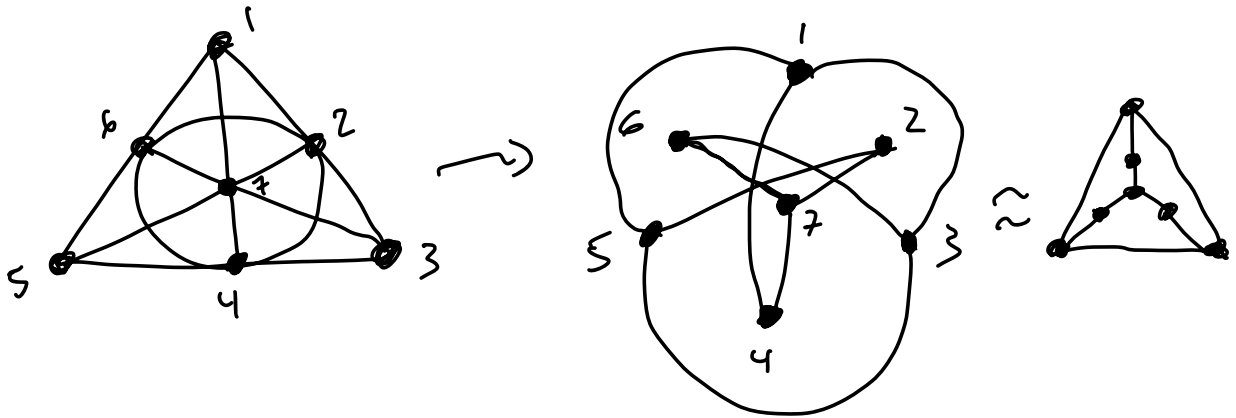
$$\begin{cases} k & \text{if } x=y & \leftarrow \text{regular deg } k \\ \lambda & \text{if } x \text{ adj. to } y \\ \mu & \text{if } x \text{ not adj. to } y \end{cases}$$

Ex.



\leftarrow Petersen Graph;
SRG w/ parameters
 $(10, 3, 0, 1)$

Construction: Given a design w/ a polarity, make a graph on X where we join $x-y$ if $x \neq y$ and $y \in \sigma(x)$



6, 4, 2 degree 2
1, 5, 7, 3 degree 3

Not regular in this example because sometimes $x \in \sigma(x)$ and other times not. But:

Lemma: If, in a polarity σ , we never (or always) have $x \in \sigma(x)$, the associated graph is strongly regular.

Pf: For never: We have $|\sigma(x) \cap \sigma(y)| = \lambda$ for all $x \neq y$ (where the design is a symmetric $2-(v, k, \lambda)$ design). So whether or not x, y are adjacent, they have λ common

neighbors. It also has degree k , so it is an (n, k, λ, μ) SRG.

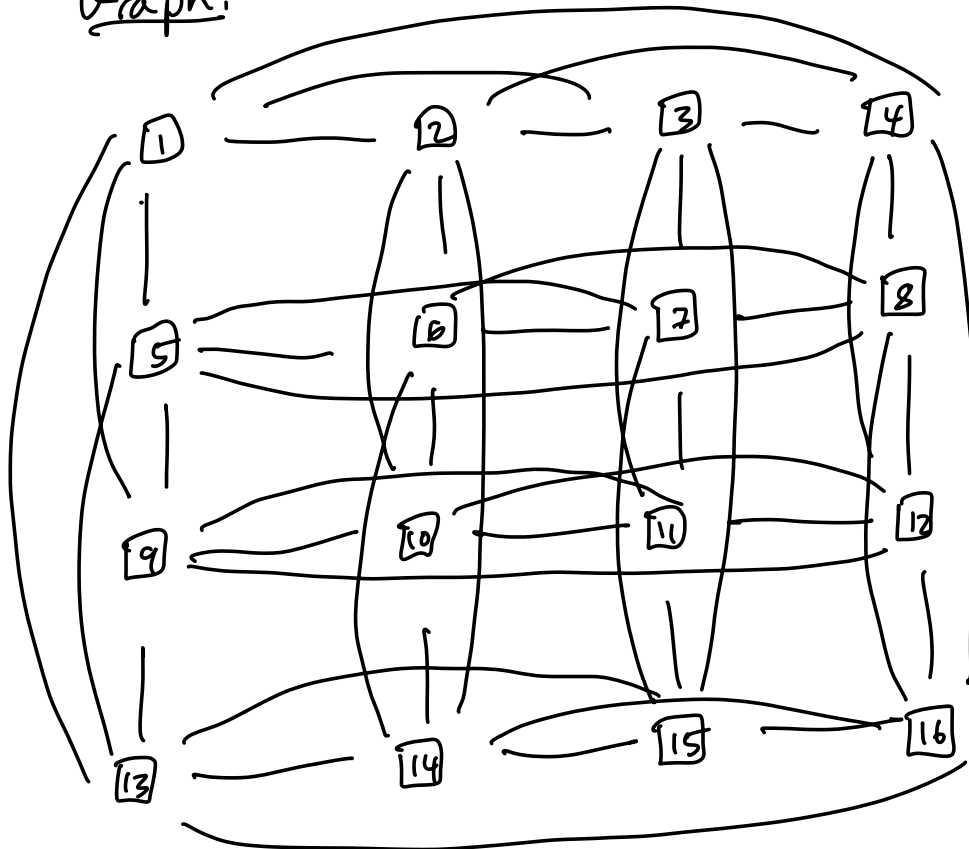
Ex:

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	16

$X = \{\text{squares at left}\}$

$B = \{2345913, 13461014, \dots\}$

Graph:



$L_2(4)$

Converse: Given an SRG (V, E) of parameters (n, k, λ, μ) , there is a symmetric $2-(n, k, \lambda)$ design and a polarity that gives rise to it.

PF: Let $X = V$ and define the blocks to be the sets of neighbors of each $x \in X$.

This gives the $2-(n, k, \lambda)$ design, and the polarity sends each x to its neighbor block. \square